

A Philosophy of Modeling

Lars Rönnbäck







No two objects can have all the same properties.

Gottfried Wilbelm Leibniz



Everything flows.

Heraclitus



The Buddhist Doctrine of Momentariness

All things are momentary.

Vasubandhu



The reality we can put into words is never reality itself.

Werner Heisenberg



Some things are more nearly certain than others.

Bertrand Russell



Our subjectivity is so completely our own.

Spike Jonze



One should always aim at being interesting, rather than exact.

Voltaire



If it quacks like a duck...

James Whitcomb Riley



Same same, but different.

Thai salesperson

When is a thing *a thing*?





When is a thing *a thing*?



- Things are important
- Their properties are important
- Their classification is ephemeral
- Values are always imprecise
- Changes can be captured through sampling
- Only identities are immutable
- We record statements about reality, and we do not record reality itself
- The certainty of true reality cannot be captured
- Opinions may differ and may be revised
- Keep the number of things that have to be agreed upon small



Transitional Modeling Lars Rönnbäck

 $[\{(,),...,(,)\},,]$

APPEARANCE

APPEARANCE SET

PRIMITIVE **STRUCTURED IMPRECISE**

[{(Identity₁, *role₁*), ..., (Identity_n, *role_n*)}, value, timepoint]

APPEARANCE TIMEPOINT

POSIT

The least we must agree upon are appearances

APPEARANCE

[{(Identity₁, *role₁*), ..., (Identity_n, *role_n*)}, value, timepoint]

APPEARANCE SET

...which lets us disagree about everything else

PRIMITIVE STRUCTURED IMPRECISE

> APPEARANCE TIMEPOINT

POSIT

[{(42, *busband*), (43, *wife*)}, married, 2004-06-19]

[{(42, *bair color*)}, gray, 2022-02-22]

[{(42, *name*)}, Lars Samuelsson, 1972-08-20]

[{(42, *bair color*)}, **brown**, 1973-02-13] [{(42, *bair color*)}, gray, 2022-02-22]

VALUES MAY CHANGE OVER TIME (APPEARANCE TIMELINE)

[{(42, *name*)}, Lars Samuelsson, 1972-08-20] [{(42, *name*)}, Lars Rönnbäck, 2005-03-30]

| DATA | [{(42, <i>name</i>)}, Lars R |
|----------|-------------------------------------|
| DATA | [{(42, <i>bair color</i>)}, gr |
| DATA | [{(4711, <i>name</i>)}, Per |
| PERIDATA | [{(42, <i>thing</i>), (4711, 6 |
| | |

RESERVED F ROLE

Rönnbäck, 2005-03-30] **ray**, 2022-02-22] **rson**, 2019-10-20] *class*)}, **active**, 1972-08-20]

RESERVED ROLE

POSIT IDENTITY $555 \leftarrow [{(42, name)}, Lars Rönnbäck, 1972-08-20]$ $556 \leftarrow [{(42, bair color)}, gray, 2022-02-22]$ $557 \leftarrow [{(42, busband), (43, wife)}, married, 2004-06-19]$

POSIT **IDENTITY**

> **555** ← [{(**4**2, *name*)}, **Lars Rönnbäck**, 1972-08-20] $556 \leftarrow [(42, bair color)], gray, 2022-02-22]$

[{(555, posit), (42, ascertains)}, 90%, 2019-10-20]

ASSERTION

ASSERTION TIMEPOINT CERTAINTY

$557 \leftarrow [{(42, busband), (43, wife)}, married, 2004-06-19]$

555 ← [{(**42**, *name*)}, **Lars Rönnbäck**, 1972-08-20] $556 \leftarrow [{(42, bair color)}, gray, 2022-02-22]$

 $557 \leftarrow [{(42, busband), (43, wife)}, married, 2004-06-19]$

[{(555, posit), (42, ascertains)}, 90%, 2019-10-20] [{(556, posit), (42, ascertains)}, -80%, 2019-10-20] [{(557, posit), (42, ascertains)}, 0%, 2019-10-21]

COMPLETE **UNCERTAINTY**

$557 \leftarrow [{(42, busband), (43, wife)}, married, 2004-06-19]$ 558 ← [{(42, *busband*), (43, *wife*)}, **married**, 2005-06-19]

CERTAINTY MAY CHANGE OVER TIME (ASSERTION TIMELINE)

[{(557, posit), (42, ascertains)}, **90%**, 2019-10-20] [{(557, posit), (43, ascertains)}, -100%, 2019-10-20] [{(557, posit), (42, ascertains)}, 0%, 2019-10-21] [{(558, posit), (42, ascertains)}, **100%**, 2019-10-21] [{(558, posit), (43, ascertains)}, **100%**, 2019-10-21]

555 ← [{(**42**, *name*)}, **Lars Rönnbäck**, 1972-08-20] 556 ← [{(42, *bair color*)}, **gray**, 2022-02-22] $557 \leftarrow [{(42, busband), (43, wife)}, married, 2004-06-19]$

[{(555, posit), (2001, job)}, success, 2019-10-20] $[\{(2001, started at)\}, 2019-10-20 10:15, 2019-10-20]$ [{(2001, *user*)}, **SQLAgent**, 2019-10-20]

METADATA

- From here we could go on and also define:
 - Constraints (and cardinality) as posits
 - Identifiers (natural keys) as posits
 - Structures (such as ensembles) as posits

• Inheritance, Trustworthiness, Consensus, Contradictions, and so on...



Schemaful Databases

How to model without classes

Lars Rönnbäck

beard color

height

soci secu numl

| | name | |
|---------------|---------------|------|
| t | hair color | |
| ial | | |
| urity nber | | RFID |
| | | |

owner, pet









DATA [{(42, name)}, Lars Rönnbäck, 1972-08-20] DATA [{(42, hair color)}, gray, 2022-02-22] DATA [{(4711, name)}, Person, 2019-10-20] PERIDATA [{(42, thing), (4711, class)}, active, 1972-08-20]

An *anchor* named PE_Person holds the identities for things of the Person class.





A static attribute named PE_NAM_Person_Name holds a reference to an identity and a primitive value.

(42, *name***)**, **Lars Rönnbäck**, 1972-08-20] [{(42, *thing*), (4711, *class*)}, **active**, 1972-08-20]



A historized attribute named PE_HAC_Person_HairColor holds a reference to an identity, a primitive value, and a time point since when it came into effect.

DATA [{(42, name)}, Lars Rönnbäck, 1972-08-20] DATA [{(42, hair color)}, gray, 2022-02-22] DATA [{(4711, name)}, Person, 2019-10-20] PERIDATA [{(42, thing), (4711, class)}, active, 1972-08-20]

A knot named COL_Color holds an enumeration of primitive values along with their own identites.



Knots are equivalent to an anchor with a single static attribute.

{{(42, *lecturer*), (911, *invoice*)}, **created**, 2019-10-20] {{(42, *busband*), (43, *wife*)}, **married**, 2004-06-19]

A static tie named PE_lecturer_IN_invoice holds references to identities in adjoined anchors.

A historized tie named PE_husband_PE_wife holds references to identities in adjoined anchors, and a time point indicating since when the relationship has been in effect.



A knotted historized tie named

references to identities in adjoined anchors, a time point indicating since when the relationship has been in effect, and the identity of an enumerated


555 ← [{(**42**, *name*)}, **Lars Rönnbäck**, 1972-08-20]

[{(555, posit), (2001, job)}, success, 2019-10-20] [{(2001, *user*)}, **SQLAgent**, 2019-10-20]

A knotted historized tie named PE_husband_PE_wife_STA_currently holds references to identities in adjoined anchors, a time point indicating since when the relationship has been in effect, and the identity of an enumerated value from a knot.

With metadata it also holds a reference to a metadata identity.

 $[\{(2001, started at)\}, 2019-10-20 10:15, 2019-10-20]$







In *concurrent-reliance-temporal* Anchor modeling an Annex table is added, in which assertions are stored.



Semantic Triples

World Wide Web Consortium

N3C°



<Bob> <is a> <Person> <Bob> <is a friend of> <Alice> <Bob> <is born on> <the 4th of July 1990> <<u>http://example.name#BS12</u>> <known as> <Bob> <<u>http://example.name#AW8</u>> <known as> <Alice>



<Bob> <is a> <Person> <Bob> <is a friend of> <Alice> <Bob> <is born on> <the 4th of July 1990> <<u>http://example.name#BS12</u>> <known as> <Bob> <<u>http://example.name#AW8</u>> <known as> <Alice>

active, 1990-07-04

active, 2014

[{(http://example.name#BS12, *known as*)}, **Bob**, 1990-07-04] [{(http://example.name#AW8, *known as*)}, Alice, 1992-10-11]



- [{(http://example.name#BS12, thing), (http://xmlns.com/foaf/0.1/Person, class)},
- [{(http://example.name#BS12, *is a friend*), (http://example.name#AW8, of)},
- [{(http://example.name#BS12, *is born on*)}, the 4th of July 1990, 1990-07-04]



<Person> <type> <Class> <is a friend of> <type> <Property> <is a friend of> <domain> <Person> <is a friend of> <range> <Person>



<Person> <type> <Class> <is a friend of> <type> <Property> <is a friend of> <domain> <Person> <is a friend of> <range> <Person>

[{(http://xmlns.com/foaf/0.1/Person, type)}, Class, 1998] [{(http://xmlns.com/foaf/0.1/Dog, *type*)}, **Class**, 1998]

[{(http://xmlns.com/foaf/0.1/Person, is a friend), (http://xmlns.com/foaf/0.1/Person, of)}, property, 1998

[{(http://xmlns.com/foaf/0.1/Dog, is a friend), (http://xmlns.com/foaf/0.1/Person, of)}, property, 1998



CONSTRAINTS

<Person> <type> <Class> <is a friend of> <type> <Property> <is a friend of> <domain> <Person> <is a friend of> <range> <Person> <is a good friend of> <subPropertyOf> <is a friend of>

- $901 \leftarrow [{(http://xmlns.com/foaf/0.1/Person, is a good friend),$ (http://xmlns.com/foaf/0.1/Person, of)}, property, 1998
- $902 \leftarrow [{(http://xmlns.com/foaf/0.1/Person, is a friend),$ (http://xmlns.com/foaf/0.1/Person, of)}, property, 1998

[{(901, constraint is), (902, constraint of)}, **sub property**, 2001]



RELATED CONSTRAINTS



Additional Resources Lars Rönnbäck

+++ Let's +++ A Let's A Again

www.anchormodeling.com



Not Secure — anchormodeling.com

Appearance is Everything

In my previous article "What needs to be agreed upon", from my series about #transitional modeling, I listed the few things that must be interpreted equally among those sharing information between them. To recall, these were identities, values, roles, and time points. If we do not agree upon these, ambiguities arise, and it is no longer certain that we are talking about the same thing. We used this to create the fundamental construct in transitional modeling; the *posit*, which is a "triple" on the

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form [{(id¹, role¹), ..., (id^N, role^N)}, value, time pc position is called a *dereferencing set*, and each (is called an appearance. An appearance consist and they will be the topic of this article.

What is interesting and different from most otl is that what the identities represent may be sul individuals exchanging information in transiti on the classifications of the things they discuss. identity 42 is thought of as a 'Living Thing' by a 'Person' by a third, a 'Customer' by a fourth, a an 'Animate Object' by a sixth a 'Transaction A

MODELING CONFLICTING, UNRELIABLE, AND VARYING INFORMATION LARS RÖNNBÄCK

Most persistent memories in which bodies of information

OURTH REVISION - 16 DECEMBER 2018

are stored can only provide a view of that information as currently is, from a single point of view, and with no respec to its reliability. This is a poor reflection of reality, because information changes over time, may have many and possibl disagreeing origins, and is far from often certain. Hereat this paper introduces a modeling technique that manage onflicting, unreliable, and varying information. In order to do so, the concept of a "single version of the truth" must be abandoned and replaced by an equivocal theory that respects e genuine nature of information. Through such, information can be seen from different and concurrent perspectives, where ach statement has been given a reliability ranging from being certain of its truth to being certain of its opposite, and when that reliability or the information itself varies over time, change are managed non-destructively, making it possible to retrieve everything as it was at any given point in time. As a result, other techniques are, among them third normal form, anchor modeling, and data vault, contained as special cases of the enceforth entitled transitional modeling

transitional, modeling, information, temporality, concurrency reliability, variability, language, vaugeness, fact

ALL CURRENT information modeling techniques result in lossy implementations, in the sense that they cannot preserve combinations of who said what, when, and how sure they were of what they were saying. Modeling such requirements explicitly using traditional techniques is complex and error-prone, and therefore

tant company Up to Change (www.uptochange.com) CHOR and TRANSITIONAL MODELING

Kingma. "Criticality of data quality as exemplified in tw disasters". In: *Information* lanagement 39.2 (2001) pp. 109–116 Krishna Kulkarni and Ja Eike Michels. "Tempora features in SQL: 2011". ACM Sigmod Record 41.3 (2012), pp. 34-43 ⁶ Tom Johnston. *Bitempora* structured way:

Rizzi. "A survey on tempor

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(IJDWM) 5.1 (2009), pp. 1-

B Liu. Uncertainty Theory

⁴ Craig W Fisher and Bruce R

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Verlag, Berlin, 2004

data warehousing". In: In

Warehousing and Minin

rarely done in practice, resulting in lost information impossible to recover (Golfarelli and Rizzi)². In fact every statement has an origin and is made with some degree of confidence (B. Liu)³, but in most modeling techniques the simplified assumption is that statement are universal, univocal, and unvarying. Riddance of such limitations should be of interest to, for example institutions subject to financial regulations that stipu late complete auditability, uncertain and complement tary clinical results within the health care domain, and conflicting information that may be strengthened or discarded over time within military, policial, or judicial applications (Fisher and Kingma)⁴. Modeled information often ends up in databases, and while database vendors have started to add rudimentary temporal capabilities (Kulkarni and Michels)⁵, these are still insufficient (Johnston)⁶ and rely on the relational model. Rather than attempting to extend the relational model, this paper introduces a new generally applicable modeling technique, that is lossless with respect to concurrency, reliability and temporality. It is simplistic in nature, yet able to model complex scenarios. It can, for example, capture the following illustrative story in a

The accused. Archie, was seen fleeing the scene of the crime by two witnesses, Donna and Charlie. Donna is h that the ac thinks he had a red beard. When the victim, Bella, later regained consciousness, she corroborated Charlie's stor at which time Donna retracted her statement. In fact, Archie had been wearing a fake red beard during the attack, but it fell off when he fled, eventually leading to his conviction through a DNA match

In the terminology introduced, *concurrency* refers to having multiple, possibly conflicting, views of the same

mporal Dimensional Modeling, Rönnbäck and Regardt, preprint 2019 rev.2

Temporal Dimensional Modeling

LARS RÖNNBÄCK OLLE REGARDT of Computer Science, Stockholm Uni

n une concept of slowly chang-ns (Kimball 2008; Ross 2013), oduced to manage change in odeling, and how well these long with the tradit 012; Rönnbäck et al. 2010;

2 LARS RÖNNBÄCK – DOI: 10.13140/RG.2.2.34381.49121/1

2014; Azvine, Cui, and Nauck 2005). Whil



www.anchormodeling.com/modeler/latest





github.com/Roenbaeck/sisula

E README.md

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sisula

sisula, short for "simple substitution language", is a language for producing text output from XML input.

The current version is built in JavaScript and should run using HTA in any Windows version from the last decade. There are no special requirements or dependencies. A legacy version using JScript in Windows Scripting Host is also available.

ETL

The ETL branch contains an SQL driven ELT framework for data warehouse automation. This framework can be used with SQL Server and is particularly useful for Anchor Modeling. There is a playlist of video tutorials on how to use it available here: https://www.youtube.com/playlist?list=PLG6-3kKEOyYIWEaEFzhcARtjqHU6zn1cH

Sisulator

The sisulator takes an XML file as input and converts this into a JSON-compatible object according to a mapping ruleset. It will then process a number of sisulets as specified in the given directive, which recieve the object as input. The sisulets are parsed and the sisula language substituted to JavaScript/JScript using regular expressions, after which the JavaScript/JScript is evaluated and the output stored.

History

sisula was introduced in Anchor Modeling in order to replace XSLT for producing text output, and a first JavaScript version of the Sisulator is built into its modeling tool. This version is derived from that work.

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2

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